

Multi-client Predicate-only Encryption for Conjunctive Equality Tests

Tim van de Kamp

Andreas Peter

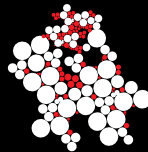
Maarten Everts

Willem Jonker



16th International Conference on Cryptology And Network Security, 2017





Monitoring over Encrypted Data

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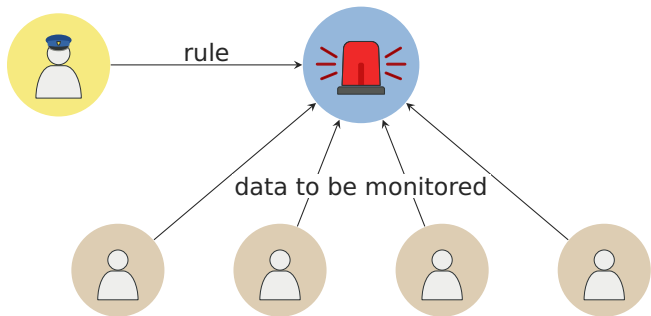
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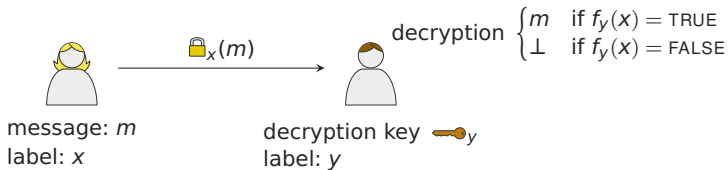


This Talk: Monitoring over Encrypted Data



Monitoring of **sensitive data** using **sensitive** monitoring **rules**.

Background: Predicate Encryption

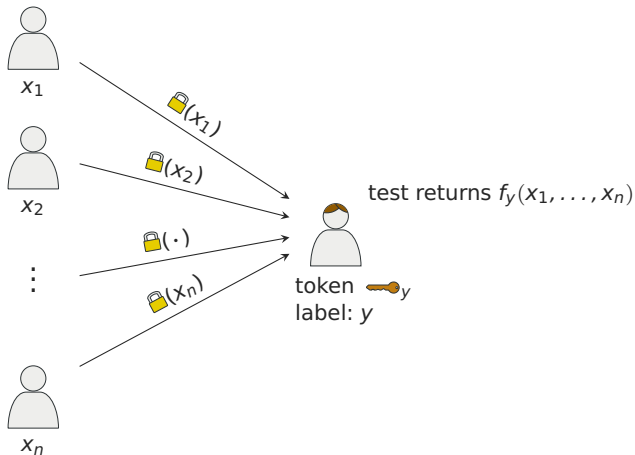


Predicate encryption for relation $R(x, y)$.

Examples

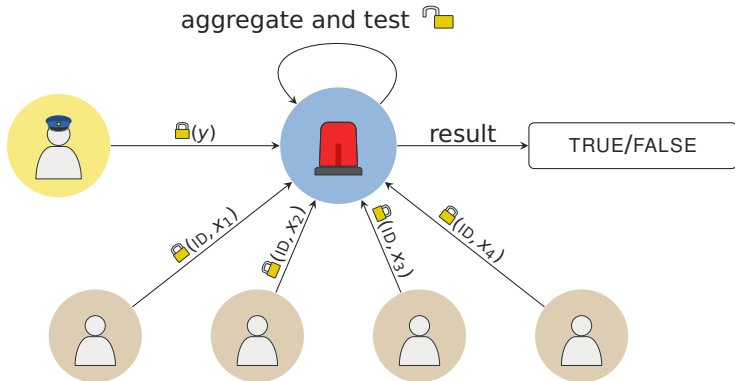
- Identity-based encryption
- Attribute-based encryption
- Hidden vector encryption
- Inner-product predicate encryption

Multi-client Predicate-only Encryption – Concept

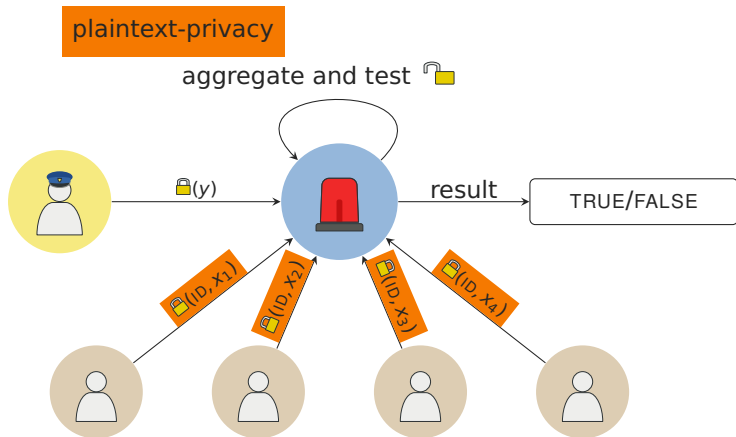


Multi-client predicate-only encryption for relation $R(x_1, \dots, x_n, y)$.

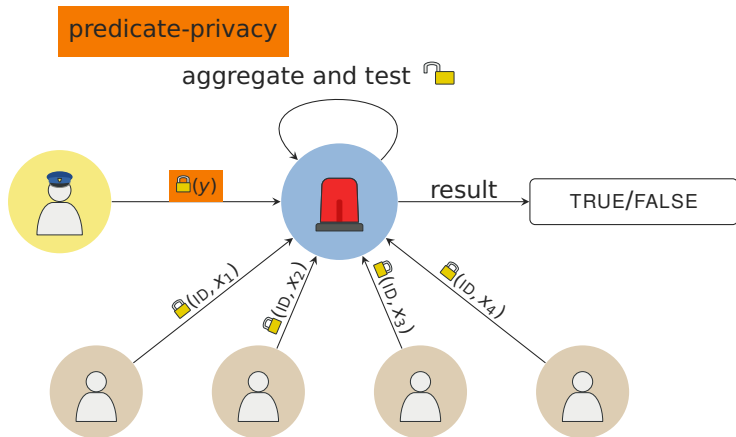
Multi-client Predicate-only Encryption – Security



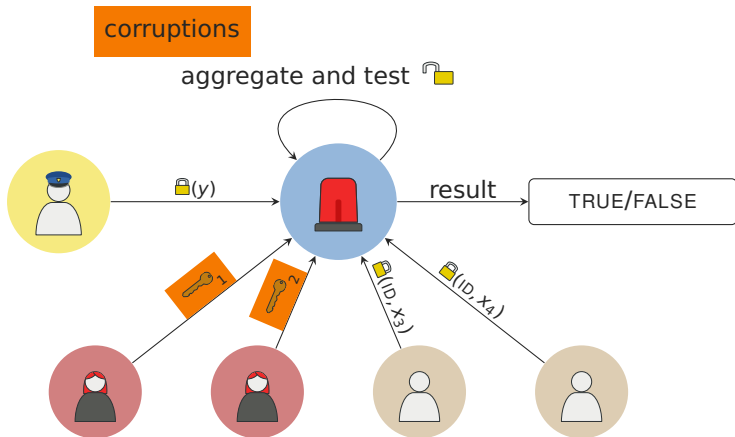
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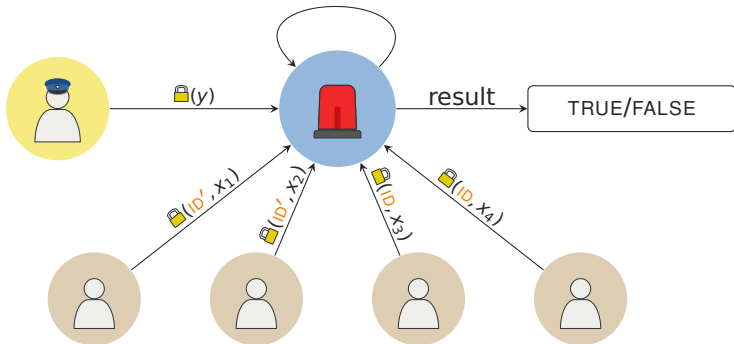
Multi-client Predicate-only Encryption – Security



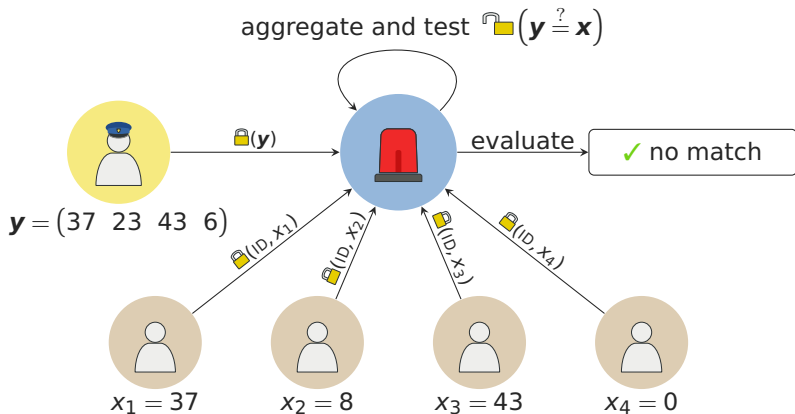
Multi-client Predicate-only Encryption – Security

mix-and-match attacks prevention

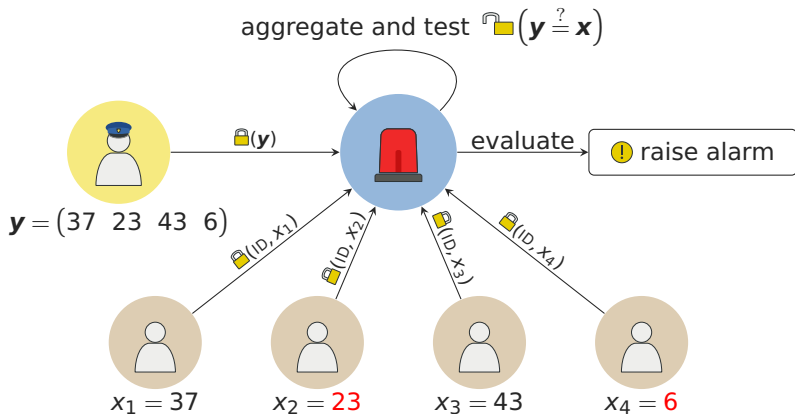
aggregate and test 



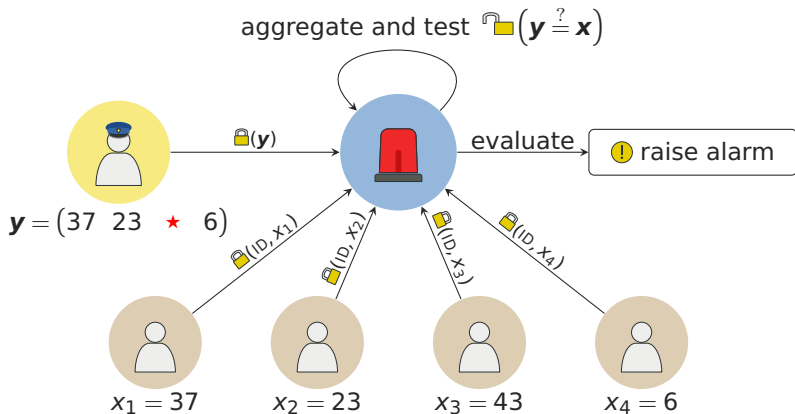
Construction: Schematic Overview



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Construction: Schematic Overview



Construction: Simplified & Highlights

Setup(1^λ)

- prime-order asymmetric pairing $e: \mathbb{G}_1 \times \mathbb{G}_2 \rightarrow \mathbb{G}_T$
- hash function $H: \{0, 1\}^* \rightarrow \mathbb{G}_1$
- $\text{usk}_i = (g_1^{\alpha_i}, \quad)$
- $\text{msk} = \left\{ (g_2^{\alpha_i}, \quad) \right\}_{i \in [n]}$

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Encrypt($\text{usk}_i, \text{ID}, x_i$)

$$\text{ct}_i = (\quad, g_1^{r_i}, g_1^{\alpha_i x_i r_i})$$

GenToken(msk, \mathbf{y})

$$\text{tk}_{\mathbf{y}} = \left(\{g_2^{u_i}, g_2^{\alpha_i y_i u_i}\}_{i \in [n]}, \quad \right)$$

Test($\text{tk}_{\mathbf{y}}, \{\text{ct}_i\}_{i \in [n]}$)

$$\prod_{i \in [n]} e(g_1^{\alpha_i x_i r_i}, g_2^{u_i}) \stackrel{?}{=} \prod_{i \in [n]} e(g_1^{r_i}, g_2^{\alpha_i y_i u_i})$$

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Encrypt($\text{usk}_i, \text{ID}, x_i$)

$$\text{ct}_i = (H(\text{ID}), g_1^{r_i}, g_1^{\alpha_i x_i r_i} H(\text{ID})^{\delta_i})$$

GenToken(msk, \mathbf{y})

$$\text{tk}_{\mathbf{y}} = \left(\{g_2^{u_i}, g_2^{\alpha_i y_i u_i}\}_{i \in [n]}, \prod_{i \in [n]} (g_2^{\delta_i})^{u_i} \right)$$

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Construction: Simplified & Highlights

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Encrypt($\text{usk}_i, \text{ID}, x_i$)

$$\text{ct}_i = (H(\text{ID}), g_1^{r_i}, g_1^{\alpha_i \pi_i(x_i) r_i} H(\text{ID})^{\delta_i})$$

GenToken(msk, \mathbf{y})

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Test($\text{tk}_{\mathbf{y}}, \{\text{ct}_i\}_{i \in [n]}$)

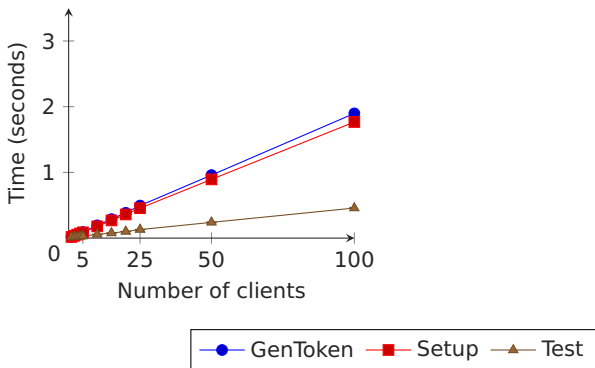
$$\prod_{i \in [n]} e(g_1^{\alpha_i \pi_i(x_i) r_i} H(\text{ID})^{\delta_i}, g_2^{u_i}) \stackrel{?}{=} \prod_{i \in [n]} e(g_1^{r_i}, g_2^{\alpha_i \pi_i(y_i) u_i}) e(H(\text{ID}), \prod_{i \in [n]} (g_2^{\delta_i})^{u_i})$$

Evaluation

Proof-of-concept implementation in Go [CRIPTIM].

MNT-159 curve

Encrypt (single client): 2.6 ms

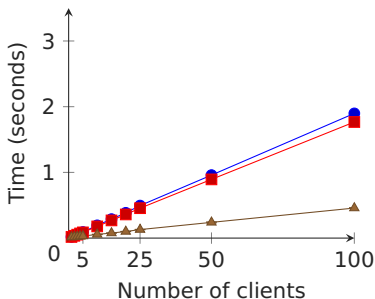


Evaluation

Proof-of-concept implementation in Go [CRIPTIM].

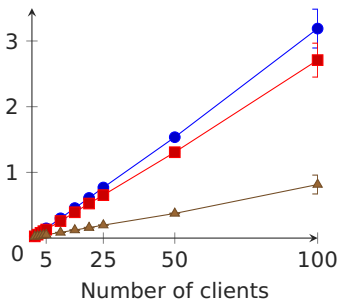
MNT-159 curve

Encrypt (single client): 2.6 ms



MNT-224 curve

Encrypt (single client): 4.4 ms



Summary

- Monitoring over encrypted data
- Defined multi-client predicate-only encryption
- **Simple** and **efficient** construction for conjunctive equality tests

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- Monitoring over encrypted data
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- **Simple** and **efficient** construction for conjunctive equality tests

Interested?

Contact: t.r.vandekamp@utwente.nl

References

[CRIPTIM] *Implementations of Private Information Sharing Schemes*. CRIPTIM consortium. URL: <https://github.com/CRIPTIM/>.



National Cyber Security Centre
Ministry of Security and Justice



Ministry of the Interior and
Kingdom Relations